

Formal Semantics for Atomic Sentences

Interpretations

- In giving the formal semantics for atomic sentences of Predicate Logic, we must introduce a number of new metavariables.
 - ‘**I**’ designates an interpretation.
 - ‘**D**’ designates a set containing at least one object (the domain of the interpretation).
 - ‘**v**’ designates a (“valuation”) function which gives the designation of constant terms, the extensions of predicates, and the truth-values of sentence letters.
 - ‘**f**’, ‘**g**’, ‘**h**’ designate function symbols.
- **I** is an ordered pair $\langle \mathbf{D}, \mathbf{v} \rangle$.

Specifying the Domain

- We may specify the domain by enumerating its members:
 - $\mathbf{D} = \{\text{Adam, Eve}\}$
 - $\mathbf{D} = \{1, 2, 3, 4, 5\}$

A domain may also be specified by description:

- $\mathbf{D} =$ the set of all positive integers
- $\mathbf{D} = \{x: x \text{ is a positive integer}\}$

More generally, we may refer to objects in the domain using the metavariables $\mathbf{u}_1, \mathbf{u}_2, \mathbf{u}_3, \dots$

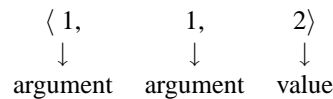
The Valuation Function

- In the following interpretation $\mathbf{I} = \{\mathbf{D}, \mathbf{v}\}$, the valuation function ‘**v**’ assigns truth-values to sentence letters.
 - $\mathbf{v}(\mathbf{P}) = \text{t}$
 - $\mathbf{v}(\mathbf{Q}) = \text{f}$
- **v** assigns members of the domain to names.
 - $\mathbf{v}(\mathbf{a}) = \text{Adam}$
 - $\mathbf{v}(\mathbf{e}) = \text{Eve}$
- **v** assigns extensions to predicates.

- $\mathbf{v(B)} = \{\langle \text{Adam} \rangle\}$
- $\mathbf{v(B)} = \{x: x \text{ is blond}\}$
- $\mathbf{v(L)} = \{\langle \text{Adam, Eve} \rangle, \langle \text{Eve, Adam} \rangle\}$
- $\mathbf{v(L)} = \{\langle x, y \rangle: x \text{ loves } y\}$

Valuations of Function Symbols

- The valuation function assigns a value to each function symbol (without parentheses).
- For a function symbol of n places, the value is a set of $n+1$ -tuples.
 - Consider the two-place addition function and a domain consisting of the positive integers.
 - $\mathbf{v(f)} = \{\langle 1, 1, 2 \rangle, \langle 1, 2, 3 \rangle, \langle 1, 3, 4 \rangle, \dots\}$.
- The first n places specify the values of the argument terms.
- The last place specifies the value of the function.



Values of Filled-in Function Symbols

- The value assigned to the function symbol allows the determination of the value of a filled-in symbol.
- Let f be a two place function.
- Then $\mathbf{v(f)}$ is a set of ordered triples.
- $\mathbf{v(f(t, u))}$ is the last member of the the ordered triple whose first two members are $\mathbf{v(t)}$ and $\mathbf{v(u)}$.
 - Let $\mathbf{v(f)}$ be the addition function as specified above.
 - Let $\mathbf{v(a)} = 1$.
 - Then $\mathbf{v(f(a,a))}$ is the last member of $\langle \mathbf{v(a)}, \mathbf{v(a)}, 2 \rangle$, which is the same as $\langle 1, 1, 2 \rangle$, and therefore is 2.

Symbols for Set Membership and Non-Membership

- It is convenient to have a symbol indicating membership in a set.
- The symbol ‘ \in ’ indicates set-membership.
 - Adam \in {Adam, Eve}
- The membership or “belongs to” symbol when struck through indicates non-membership.
 - Cid \notin {Adam, Eve}

Truth-Definition for Atomic Sentences

- Let \mathbf{P} be an n-place predicate.
- Let $\mathbf{t}_1, \mathbf{t}_2, \dots, \mathbf{t}_n$ be constant terms.
- Let $\mathbf{P}\mathbf{t}_1\mathbf{t}_2 \dots \mathbf{t}_n$ be an atomic sentence of Predicate Logic consisting of the n-place predicate \mathbf{P} followed by n constant terms $\mathbf{t}_1, \mathbf{t}_2, \dots, \mathbf{t}_n$, in that order.

Truth-Definition for Atomic Sentences of Predicate Logic

$\mathbf{v}(\mathbf{P}\mathbf{t}_1\mathbf{t}_2 \dots \mathbf{t}_n) = \mathbf{t} \text{ if } \langle \mathbf{v}(\mathbf{t}_1), \mathbf{v}(\mathbf{t}_2) \dots \mathbf{v}(\mathbf{t}_n) \rangle \in \mathbf{v}(\mathbf{P})$

$\mathbf{v}(\mathbf{P}\mathbf{t}_1\mathbf{t}_2 \dots \mathbf{t}_n) = \mathbf{f} \text{ otherwise}$
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- Compare this definition with the informal definition given earlier.
- An atomic sentence with an n-place predicate is true in an interpretation if the n-tuple consisting of the referents of the arguments (from left to right) is a member of the extension of the predicate; the sentence is false otherwise.

An Example

- Let $\mathbf{I} = \{\mathbf{D}, \mathbf{v}\}$.
- $\mathbf{D} = \{\text{Adam, Eve}\}$
- $\mathbf{v}(\mathbf{a}) = \text{Adam}$
- $\mathbf{v}(\mathbf{e}) = \text{Eve}$
- $\mathbf{v}(\mathbf{L}) = \{\langle \text{Adam, Eve} \rangle, \langle \text{Eve, Adam} \rangle\}$
- So, $\langle \text{Adam, Eve} \rangle \in \mathbf{v}(\mathbf{L})$
- So, $\langle \mathbf{v}(\mathbf{a}), \mathbf{v}(\mathbf{e}) \rangle \in \mathbf{v}(\mathbf{L})$
- So, ‘Lae’ is true in \mathbf{I} .

Another Example

- Let $I = \{\mathbf{D}, \mathbf{v}\}$.
- \mathbf{D} = the set of all positive integers.
- $\mathbf{v}(o) = 1$
- $\mathbf{v}(t) = 2$
- $\mathbf{v}(f) = \{\langle 1, 1, 2 \rangle, \langle 1, 2, 3 \rangle, \dots\}$ [the addition function]
- $\mathbf{v}(G) = \{\langle x, y \rangle \mid x > y\}$
- So, $\mathbf{v}(f(o, o)) = 2$ [$1 + 1 = 2$]
- So, $\langle 2, 2 \rangle \notin \mathbf{v}(G)$
- So, $\langle \mathbf{v}(t), \mathbf{v}(f(o, o)) \rangle \notin \mathbf{v}(G)$
- So, ' $Gf(o, o)$ ' is false in I .